



SMM Protection in EDK II

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Agenda





- More Protection
 - SMM Memory Protection
 - -CommBuffer Enforcement
 - -ASLR in SMM
 - Guard Page
 - -Reduce SMI Handler
- Summary / Call to Action

What is SMM and SMI?



- System Management Mode (SMM)
 - Is a special CPU operating mode.
 - Is inside of a special SMM memory (SMRAM)
 - Access the whole system memory and IO, including OS memory and hypervisor memory.
 - Is invoked through a System Management Interrupt (SMI)
 - Has software executive (SMI handler) to perform operation based upon different SMI.

Known SMM Attacks



| SMM Attack | Description | Example |
|---------------------------------|---|----------|
| SMRAM is unlocked | An attacker can set register to unlock SMRAM, and override SMRAM. | A.1, A.2 |
| Cache Poisoning | An attacker can set CPU cache to override SMRAM. | A.3, A.4 |
| SMRAM remap | An attacker can control chipset register to remap a normal system memory to SMRAM. | A.5 |
| Branch Outside of SMRAM | SMM code calls outside of SMRAM, which is controlled by the attacker. | A.6, A.7 |
| SMM Communication Buffer Attack | SMM code uses SMM communication buffer to exchange information with non-SMM agent. The attacker can give a malicious communication buffer to SMM, and the SMM may write SMRAM or Virtual Machine Monitor (VMM). | A.8, A.9 |

Known Mitigation



| SMM Attack | Mitigation |
|-----------------------------|--|
| SMRAM is unlocked | 1) Lock SMRAM at PI SmmReadyToLock. |
| Cache Poisoning | 1) Enable SMM Range Register. |
| SMRAM remap | 1) Lock Remap register. |
| Branch Outside of SMRAM | Enable Smm_Code_Access Register. Setup Non-Executable (NX) paging outside of SMM. |
| SMM Communication Attack | Check SMM Communication Buffer. Check MemoyMapped IO (MMIO) bar access. |

New methods may be discovered



Current SMRAM Layout

- Every page in SMRAM is read/write
- Every page in SMRAM is executable

| MSEG | |
|----------------------|--|
| PiSmmCore (PE/COFF) | |
| SMM Driver (PE/COFF) | |
| PiSmmCpu (PE/COFF) | |
| SMM Save State | |
| SMM Stack | |
| SMM IDT/GDT | |
| SMM Page Table | |
| SMM Driver (PE/COFF) | |
| Other Heap Data | |
| ••• | |
| SMM S3 Resume State | |

SPA -

MSEG

PiSmmCore (PE/COFF)

SMM Driver (PE/COFF)

PiSmmCpu (PE/COFF)

SMM Save State

SMM Stack

SMM IDT/GDT

SMM Page Table

SMM Driver (PE/COFF)

Other Heap Data

•••

SMM S3 Resume State

SMM Memory Protection

PE Header

PE Code

PE Data

CPU m Save State

CPU 2m - 1 SMI Entry

CPU m-1 Save State

CPU 2m-2 SMI Entry

CPU 1 Save State

CPU m SMI Entry

CPU O Save State

CPU m-1 SMI Entry

pad

CPU 1 SMI Entry

pad

CPU 0 SMI Entry

RO

XD

TSEG

Using static page table

 Set NonExecutable (NX) for data

- Set ReadOnly(RO) for code
- Protect page itself
- SMM driver can protect its own critical data in ReadOnly(RO) memory

PE Code PE Data Page Table SMM EntryPoint Save State IDT/GDT Data (Stack/Heap) **RO-Data**



- Prevents code injection
- Protects critical data (read-only)

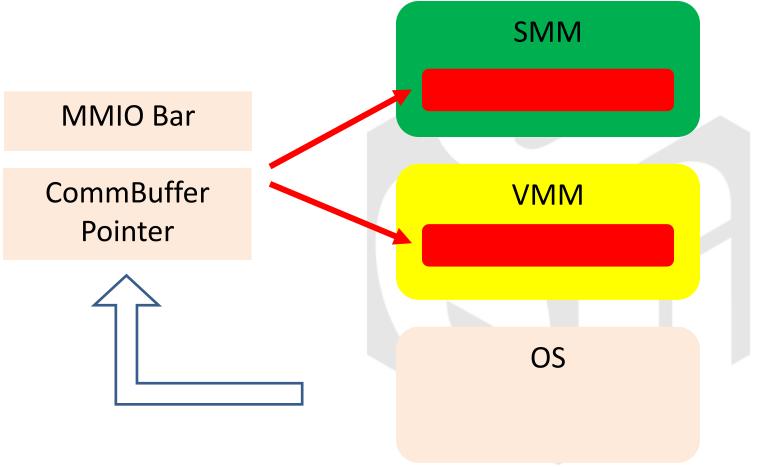
- Limitations
 - Return-oriented programming (ROP) attack.
 - Size overhead
 - PE image: 6K * SmmImageCount (average)
 - Static Page Table: 2M (1G paging for 48bit)



CommBuffer Enforcement

SMM CommBuffer Attack





Current CommBuffer Check

 SMI handler MUST check SMM communication buffer content by writing code like below:

```
if (!SmmIsBufferOutsideSmmValid ((UINTN)CommBuffer, TempCommBufferSize)) {
   DEBUG ((EFI_D_ERROR, "SmmVariableHandler: SMM communication buffer in
   SMRAM or overflow!\n"));
   return EFI_SUCCESS;
}
```

• But if the check is missing, there is no way to detect such problem.

CommBuffer Enforcement

BootCode BootData Valid MMIO **MMIO SMRAM** Policy Enforcement Reserved Valid **SMM ACPINvs** Page Table **SmmComm** RuntimeCode Buffer RuntimeData **ACPI Reclaim BootCode BootData** LoaderCode **Not-Present** LoaderData

CommBuffer Enforcement

- Resist comm buffer attack even the CommBuffer check is missing in SMM driver
- Protects hypervisors
- Limitation
 - This enforcement is not applied to hotplug memory, which is still read/write.
 - MMIO region is mapped. SMI handler need make sure MMIO bar point to a valid region.



Address Space Layout Randomization (ASLR) in SMM

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Current SMM Layout

- The layout is fixed.
- This attack may find out a sequence of instruction existed in code region (gadgets) and execute. (ROP)
- This ROP attack can bypass NX/RO protection.

| MSEG | |
|----------------------|--|
| PiSmmCore (PE/COFF) | |
| SMM Driver (PE/COFF) | |
| PiSmmCpu (PE/COFF) | |
| SMM Save State | |
| SMM Stack | |
| SMM IDT/GDT | |
| SMM Page Table | |
| SMM Driver (PE/COFF) | |
| Other Heap Data | |
| | |
| SMM S3 Resume State | |

Image Shuffle



| Image A |
|---------|
| Image C |
| Image B |
| Image D |
| |

| Image C |
|---------|
| Image B |
| Image D |
| Image A |
| |

| Image D |
|---------|
| Image B |
| Image A |
| Image C |
| |

1st Boot 2nd Boot 3rd Boot

As such, it makes difficult for attacker to locate gadgets for ROP attack.

Heap Shift

It makes difficult for attacker to locate gadget for ROP attack.

MAX - Random SmmCore Random MAX - Random SMM Heap MAX-Random Stack Random MAX-Random PageTable Random Random



Fixed Length

2(Entropy)

2(Entropy)

ASLR in SMM

 Make Buffer Overflow/ROP attack harder, because the memory layout is changed in each boot.

Limitations

- SMM is a resource constrained environment.
 Entropy for Heap Shift might be not so big.
- Information leakage in SMM (LoadedImageProtocol)



Guard Page

Current Memory Allocation

- Page overflow cannot be detected
- Pool overflow can only be detected when memory is freed, because of POOL_TAIL signature check at FreePool()

Allocated Page

POOL_HEAD

Allocated Pool

POOL_TAIL

POOL_HEAD

Allocated Pool

POOL TAIL

Allocated Page

New Page Allocation



Guard Page (Not Present)

> Allocated Pages

Guard Page (Not Present) One Allocation for AllocatePages()

2 guard pages (8K)

New pool allocation



Guard Page (Not Present)

POOL_HEAD (phd0)

Allocated Pool

POOL_TAIL (ptal)

Guard Page (Not Present) One Allocation for AllocatePool()

2 guard pages (8K) + 4K page alignment

Guard Page

- Catch page overflows when they happen
- Catch pool overflows when they happen
- Limitation
 - Memory size overhead
 - Additional 8K for each page allocation.
 - Additional 8K+4K alignment for each pool allocation.
 - It might need above 128M SMRAM.
 - A debug feature, because of size overhead.



Reduce SMI Handler

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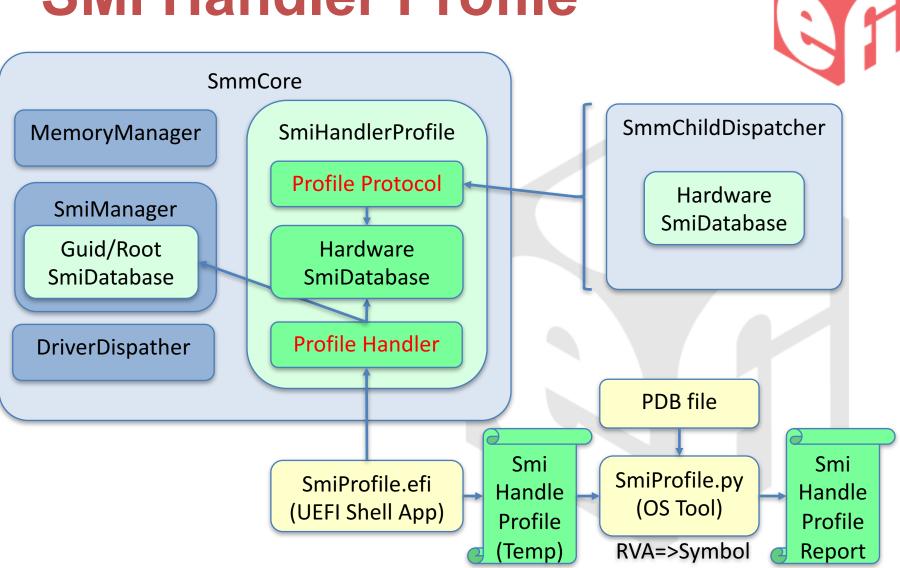
SMI Handlers



SMI Handler == Attack Surface

- Question:
 - How many SMI handlers in the BIOS?
 - How many Root SMI handlers, GUID handlers, software SMI handlers,?

SMI Handler Profile



SMI Handler Profile

- Developer can check if the SMI handler is necessary
- Test engineer can use it for validation
- Limitation
 - Only used as a debug feature (info leakage)
 - The profile only shows info, which requires further analysis

Summary

- SMM is a target due to high execution privilege
- There are known SMM attacks and mitigations
- Developers can do more to protect SMM
 - SMM Memory Protection
 - CommBuffer Enforcement
 - ASLR in SMM
 - Guard Page
 - Reduce Number of SMI Handlers

Call To Action



- Adopt "SMM Memory Protection" and "CommBuffer enforcement" to harden the platform. [P.1][P.2]
- Use "SMI handler profile" to audit the SMI handlers. [P.3]
- Evaluate "ASLR in SMM" and resolve information leakage. [P.4]
- Use "GuardPage" to validate buffer overflow.
 [P.4]

Acknowledgement



 Some content of the material is discussed with UEFI BIOS and security experts

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Reference



Attacks

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Protection:

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 (https://github.com/tianocore-docs/Docs/raw/master/White_Papers/A_Tour_Beyond_BIOS_Secure_SM_M_Communication.pdf)
- [P.2] A_Tour_Beyond_BIOS_Memory_Protection_in_UEFI
 (https://www.gitbook.com/book/edk2-docs/a-tour-beyond-bios-memory-protection-in-uefi-bios/details)
- [P.3] SMI Handler Profile Feature
 (https://github.com/tianocore/tianocore.github.io/wiki/SMI-handler-profile-feature)
- [P.4]
 A_Tour_Beyond_BIOS_Securiy_Enhancement_to_Mitigate_Buffer_Ove rflow_in_UEFI

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